

A PAL Program of the Blackboard Evaluator

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May 16, 1968

Attached is a listing of a PAL program that simulates the left-hand blackboard evaluator of 6.231. It includes the following:

- the applicative subset
- tuples
- assignment
- recursion of single functions
- simultaneous definition
- labels, in PAL's complete generality
- val and res

It does not include

- simultaneous assignment
- simultaneous recursion
- the operator jj

Adding simultaneous assignment is easy; adding jj is not too hard; and adding simultaneous recursion is much harder.

The attached listing, with pages numbered 1 to 21, contains five files, like this.

AA	page 1
BB	6
T7	8
CC	12
/LB02/	18

They are organized as follows:

AA PAL

This file contains definitions of the data structures needed, along with routines to operate on them. The order, roughly, is

- debugging
- control and stack
- environment
- memory
- codes used in control and stack
- cycle counter

dump
initialization routines

In general, the actual structure of the various data bases is known only to the routines in AA.

BB PAL

This file is best ignored. It contains a routine for printing values produced by the interpreter, and it contains various debug things.

T7 PAL

This file contains sample input to the evaluator. The comment shows the PAL program being simulated, and the rest of the file contains the definition of a tuple, D, which is the control structure.

CC PAL

This file does the work. After the definition of the function Apply-Closure, the rest of the file is imperative. The actual evaluator is on pages 13 to the middle of 16. Each label BRn processes a single control item, as indicated on page 3.

The code on page 16 labelled MAIN accepts a directive from the console and then carries it out. The possible directives are at the bottom of page 16.

/LB02/ PAL

This library file is used instead of /LB00/. It is shown here for completeness, but is not part of the evaluator.

```
// Simulate the blackboard machine executing PAL.  
  
// Arthur Evans, Jr. 24 April 68  
// This file was last modified on 05/12/68 at 10:57 by Evans.  
  
// Some definitions for debugging.  
  
def ERRCREXIT = nil // updated later to MAIN  
and CLDjj = nil // hold jj evaluated in Report  
  
def Report t = // error detected by the simulator  
    Write 'ERRCR: '; Write t; Write '*n';  
    CLDjj := jj; // save 'j' for calling EDebug  
    gotc ERRCREXIT  
  
// Define the stack and the control, and the relevant selectors.  
  
def S = nil // the stack  
and C = nil // the control  
  
// The stack and the control have the same format: Each is a  
// 3-tuple, whose third element is a 3-tuple, etc. The first  
// element of the tuple is the type of the item, and the second  
// element is relevant data (or nil).  
  
// The following three functions are selectors for the stack  
// and the control:  
  
def CODE ST = ST 1 // select code  
and DATA ST = ST 2 // select data  
and LINK ST = ST 3 // select link  
  
// Routines for the stack and control. CLDTCP is for debugging.  
  
def CLDTCP = nil  
  
def Pop ST = // pop C or S and return old top element  
    Null ST -> Report 'S or C unexpectedly empty' ! dummy;  
    CLDTCE := CODE ST, DATA ST;  
    ST := LINK ST; // pop the thing  
    $ CLDTCE  
  
and Push (x, ST) = // push x into C or S  
    ST := x 1, x 2, $ ST  
  
def Stack x = // push x into S  
    Push(x, S)
```

```

// Definitions for the ENVIRONMENT.

def E = nil           // the environment
and CurE = C          // current environment level
and LastE = C          // last environment level used
and ENVstack = nil    // stack of environment layers

// An environment level is a 3-tuple, like this:
//   1 name (as a string)
//   2 value
//   3 index in E of next level to look in

def LookupinE Name = // look up Name in current environment
  f CurE
  where rec f x =
    x = C -> Report('Can\'t find ', Name, ' in environment.')
    ! E x 1 = Name -> E x 2
    ! f(E x 3)

def EnterE(Name, Store, Link) =
  E := E aug (Name, $ Store, $ Link);
  LastE := LastE + 1;
  ENVstack := $ CurE, $ ENVstack; // remember current environment
  CurE := LastE

and PopENV () = // pop the environment stack
  CurE, ENVstack := ENVstack

and InitE () = // initialize the environment to empty
  E, CurE, LastE, ENVstack := nil, C, 0, nil

def RES = 'RES' // identifier for VAL and RFS

```

```

// Definitions for the MEMORY.

def GUESSmark = 'GUESS' // mark for guesses, for Y
and GUESScount = 0 // make all guess distinct

def M = 0, (GUESSmark, 0), nil // the memory
and LastM = C // the last cell used in the memory

def Store v = // the routine for curly S
  LastM := LastM + 1;
  M := ↓ LastM, $ v, ↓ M;
  $ LastM

and Contents c = // the routine for curly C
  let rec f x = c = x 1 -> x 2 ! f(x 3)
  in
  let t = f M // look up c in the memory
  in
  CODE t = GUESSmark // did we find a guess?

```

' -> Report 'GUESS found in memory'
 ! t // all OK, so return the value found

def Update(c, v) = // update cell c with value v
 M := \$c, \$v, \$M

and InitM () = // initialize the memory
 M, LastM, GUESScount := (0, (GUESSmark, 0), nil), 0, 0

def GUESS () = // return the next guess
 GUESScount := !GUESScount + 1;
 GUESSmark, !GUESScount

// Codes used in the CONTROL of the machine:

// name	ccode	data
def mENV	= 1	// n (index in E)
and mVAB	= 2	// print name, as a string
and mCCN	= 3	// value
and mEASIC	= 4	// value
and mLAMIA	= 5	// (body, (bv1, bv2, ...))
and mGAMMA	= 6	
and mBETA	= 7	
and mDEITA	= 8	// n, the subscript in D
and mTAU	= 9	// n, the order of the tuple to make
and mCurlyS	= 10	
and mCurlyC	= 11	
and mUPDATE	= 12	
and mSEMICCICN	= 13	
and mRVGAMMA	= 14	// n, number of operands (1 or 2)
and mY	= 15	
and mEQUALCCICN	= 16	// reverse update, for recursion
and mEIGDEITA	= 17	// (body, (L1,D1), (L2,D2), ...)
and mGC	= 18	// (body, dump)
and mVAL	= 19	
and mRES	= 20	

// Codes used in the STACK of the machine.

def ENVmark	= '	ENV'	// n, the index in E
and LVmark	= '	IV'	// n, memory location
and RVmark	= '	RV'	// value
and EASICmark	= '	EASIC'	// value
and LAMIAmark	= '	LAMIA'	// (body, (bv1, bv2, ...), env)
and TUPLEmark	= '	TUPLE'	// n, memory location
and DUMMYmark	= '	DUMMY'	
and FPmark	= '	FE'	// (n, line) (n is index in D)

def FALcycles = 0

// Stuff to count cycles of the evaluator.

```
def IntCYCLES = 0 // total count of interpretation cycles
and MaxCYCLES = 100 // max cycles permitted per evaluation
and StartCYCLES = 0 // cycle count at beginning of an execution
and CurCYCLES = 0 // cycles since last execution of InitCY

def CountCYCLES () = // count interpretation cycles
    IntCYCLES := IntCYCLES + 1;
    IntCYCLES > MaxCYCLES
    -> Report 'Interpretation cycle limit exceeded.'
    ! dummy

def InitCY () = // initialize stuff for cycles
    StartCYCLES := IntCYCLES;
    CurCYCLES := IntCYCLES

and MoreCY () =
    CurCYCLES := IntCYCLES

def LastDELTIA = 0 // last delta encountered, for debugging

// Definitions for the DUMP.

def Dump = nil // The dump, itself.
and DumpCount = 1 // The next dump layer to be assigned.

// The dump is a k-tuple, where k is the number of layers currently
// in use. Each layer is a 4-tuple:
//      C, S, ENVstack, CURE

def AddDump () = // add a layer to the dump
    Dump := Dump aug (SC, BS, ENVstack, CURE);
    DumpCount := DumpCount + 1

and RestoreDump n = // restore things as of dump layer n
    let Du = Dump n
    in
    C := Du 1;
    S := Du 2;
    ENVstack := Du 3;
    CURE := Du 4

def InitDu () = // initialize the dump
    Dump := nil;
    DumpCount := 1

// Initialization routines for stack, control and everything.

def InitC () = // initialize the control
    C := #DELTIA, 1, nil;
    LastDELTIA := 0

and InitS () = S := nil
```

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```
def InitAll () =  
    InitC nil; // control  
    InitS nil; // the stack  
    InitE nil; // the environment  
    InitM nil; // the memory  
    InitDu nil; // the dump  
    InitCY nil; // the cycle counter  
    CLDjj := nil // jj as in Report  
  
def EMPTY = '' // the empty string  
and NEWLINE = '*n'
```

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```
// Part 2 of the Blackcard Evaluator: Debug and print routines.  
// Last modified on 05/11/68 at 21:17 by Evans.  
  
def  
  rec IF v =  
    CCDE v = TUELEMARK  
    -> ( let t = Contents(DATA v)  
        in  
        Write '(';  
        F(1, Order t);  
        Write ')'  
        where rec P(i, n) =  
          n = 0 -> Write nil  
          ! i > n -> dummy  
          ! ( IF(Contents(t i)); P(i+1, n) )  
        )  
        ! Write (' ', DATA v, ' '))  
within  
  IntErint x =  
    write NEWLINE;  
    IF( Contents(DATA x) );  
    write '*n*n'  
  
def BasicPrint x = // the routine called by code  
  IntErint x;  
  LVmark, C  
  
def CLDLocp = nil // This item is updated later to f Locp.  
def T, TA, TE, TC = nil, nil, nil, nil // some temps  
and SA, SE = nil, nil // some more  
  
def CLDSYSTEMERRCR = $ SYSTEMERRCR  
  
def CLDSYS () = // set SYSTEMERRCR as it originally was  
  SYSTEMERRCR := CLDSYSTEMERROR  
  
def SYSTEMERRCRVALUE = nil // value returned by SYSTEMERRCR  
  
def MYSYSTEMERRCR t = // my version of SYSTEMERRCR  
  let j = jj // to call EDebug  
  in  
  let icrrE = $ SYSTEMERRCR  
  in  
  SYSTEMERRCRVALUE := t;  
  SYSTEMERRCR := (11 x.x);  
  EDebug j;
```

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```
SYSTEMERROR := 1000;
$ SYSTEMERRORVALUE

def MYSYS () =
    SYSTEMERROR := MYSYSTEMERROR

def Test = nil // updated later to a little test function

def Branch = nil // Updated later to a label tuple.
```

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```

// Generalized labels.
// Last modified on 05/11/68 at 15:14 by Evans.

// The FAL program:

let a, k, d = -1, nil, nil
in
let f x =
  l: a := a + 1;
    Print a;
  #: x < 0 -> (k := d)
  ! x = 0 -> (k := 1)
  ! (x := x - 3)
in
d := n; f a; qctc k;
n: a

// Operators:

def ADD x y = x + y
and SUB x y = x - y
and NEG x = - x
and LES x y = x < y
and EQU x y = x = y

def NIL = nil

// The control structure:

def C =
  // C 1: GAMMA (LAMBDA 2 (a k d)) S T3 S RVGAMMA1 NEG 1 S nil S nil
  (mGAMMA, NIL,
   (mLAMDA, (2, ('a', 'k', 'd'))),
   (mCurlyS, NIL,
    (mTAU, 3,
     (mCurlyS, NIL,
      (mRVGAMMA, 1,
       (mBASIC, NEG,
        (mCON, 1,
         (mCurlyS, NIL,
          (mCON, NIL,
           (mCurlyS, NIL,
            (mCON, NIL,
             nil ))))))))) )
  ,
  // C 2: GAMMA (LAMBDA 3 f) S (LAMBDA 8 x)

```

(mGAMMA, NIL,
(mLAMDA, (3, 't')),
(mCurlyS, NIL,
(mLAMDA, (8, 'x'),
nil))))

,

// D 3: (BIGDELTA 4 (n 7))

(mBIGDELTA, (4, ('n', 7)),
nil)

,

// D 4: DS ; := d C n

(mDELTA, 5,
(mSEMICCICK, NIL,
(mULATE, NIL,
(mVAB, 'd',
(mCurlyC, NIL,
(mVAR, 'n',
nil)))))

,

// D 5: D6 ; GAMMA C f a

(mDELTA, 6,
(mSEMICCICK, NIL,
(mGAMMA, NIL,
(mCurlyC, NIL,
(mVAB, 'f',
(mVAB, 'a',
nil)))))

,

// D 6: D7 ; GO C k

(mDELTA, 7,
(mSEMICCICK, NIL,
(mGO, NIL,
(mCurlyC, NIL,
(mVAB, 'k',
nil)))))

,

// D 7: a

(mVAB, 'a',
nil)

,

// D 8: (FIGDELTA 9 (l 9) (m 11))

(mBIGDELTA, (9, ('l', 9), ('m', 11)),
nil)

// D 9: D10 := a RVGAMMA + C d 1

(mDELTA, 10,
(mSEMICICK, NIL,
(mUPDATE, NIL,
(mVAR, 'a',
(mRVGAMMA, 2,
(mBASIC, ADD,
(mCurlyC, NIL,
(mVAR, 'a',
(mCON, 1,
nil))))))))

,

// D 10: D11 := PRINT a

(mDELTA, 11,
(mSEMICICK, NIL,
(mGAMMA, NIL,
(mBASIC, BasicPrint,
(mVAR, 'a',
nil)))))

,

// D 11: D12 D13 BETA RVGAMMA LES C x C

(mDELTA, 12,
(mDEITA, 13,
(mBETA, NIL,
(mRVGAMMA, 2,
(mBASIC, LES,
(mCurlyC, NIL,
(mVAR, 'x',
(mCON, 0,
nil))))))

,

// D 12: := k C d

(mDATE, NIL,
(mDIA, 'k',
(mCurlyC, NIL,
(mVAR, 'd',
nil))))

,

// D 13: D14 D15 BETA RVGAMMA EQU C x C

(mDEITA, 14,
(mDELTA, 15,
(mBETA, NIL,
(mRVGAMMA, 2,
(mBASIC, EQU,
(mCurlyC, NIL,

(mVAR, 'x',
(mCON, 0,
nil))))))

// D 14: := K C 1

(mUPDATE, NIL,
(mVAR, 'k',
(mCurlyC, NIL,
(mVAR, '1',
nil))))

// D 15: := X RVGAMMA - C X 3

(mUPDATE, NIL,
(mVAR, 'x',
(mRVGAMMA, 2,
(mBASIC, SUE,
(mCurlyC, NIL,
(mVAR, 'x',
(mCON, 3,
nil))))))

CC PAL 05/12/68 1212.0 2258

```
// Part 3 of the Blackcard evaluator.  
  
// This section was last modified on 05/12/68 at 11:33 by Evans.  
  
let ApplyClosure(Rator, Rand) = // apply a lambda-closure  
    let Epart = & Rator 3 // environment in which closure formed  
    in  
        Apply(Rator 2, DATA Rand) // call the routine to do the work  
        where rec Apply(F, X) = // apply F to X  
            Istuple F // check for multiple tv-part  
            -> ( let y = Contents X // the tuple rand  
                in  
                    CCEE y = TUPLEmark  
                    -> ( let n = Order F // rator  
                        and d = Contents(DATA y) // rand  
                        in  
                            n = Order d  
                            -> ( Ap 1 // apply tuple  
                                where rec Ap k =  
                                    k > n -> dummy  
                                    ! (Apply(F k, d k); Ap(k+1))  
                                )  
                                ! Report 'Conformality failure'  
                            )  
                            ! Report 'Conformality failure.'  
                        )  
                    ) // single tv-part  
                )  
                EnterE (F, X, Epart);  
                Stack( ENVmark, $CURE );  
                Push( (ENV, $CURE), C );  
                Epart := CURE  
            )  
        in  
  
// *****  
// EXECUTION STARTS HERE!  
  
Test :=  
( f  
  where f n =  
    Lccp := T1;  
    goto CIDLccp;  
  T1: n := n - 1;  
    n > 0 -> goto CLELccp !  
    Lccp := CLELccp;  
    S
```

```

};

CLEIcop := Icop;

Branch := ER1, ER2, ER3, ER4, ER5, ER6, ER7, ER8, ER9, ER10,
         ER11, ER12, ER13, ER14, ER15, ER16, ER17, ER18, ER19, ER20;

ERROREXIT := MAIN; // exit from the routine Report
MYSYS nil; // establish my SYSTEMSRCR

goto MAIN;

// Here is the MAIN PROCESSING LOOP of the evaluator.

Loop:   Null C -> goto Done ! dummy; // All done, so quit.
        CcountCYCLES nil; // count interpretation cycles
        T := Pop C; // get the next control item
        goto Branch(CCLE T); // branch on its code

ER1:    // environment marker
        TA := Pop S; // the useful value in the stack
        TE := Top S; // the environment marker in the stack
        // Now we have a validity check.
        & (CCDE TE = ENVmark) // TE must be an environment marker
        & (DATA TE = DATA T) // it must be same env as in control
        & (DATA TE = CURE) // and same env as current
        & (CCDE TA = LVmark) // value must be an LV
        -> Stack(LVmark, DATA TA) // all OK
        ! Report 'Improper environment marker.';
        PopENV nil; // set CURE
        goto Loop;

ER2:    // variable
        Stack(LVmark, LookupinE(DATA T));
        goto Icop;

ER3:    // constant
        Stack(FVmark, DATA T);
        goto Icop;

ER4:    // basic operator
        Stack(BASICmark, DATA T);
        goto Icop;

ER5:    // lambda expression
        Stack(LAMBDAmark, DATA T aeq $ CURE );
        goto Icop;

ER6:    // gamma
        T := Pop S; // the rator
        TA := Pop S; // the rand
        CCDE T = LAMBDAmark // check the type of the rator
        -> ( // apply a lambda closure
            ApplyClosure(DATA T, TA); // apply the closure
            Push( (WDELTA, DATA T), C)

```

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```

)
!
CCDE T = TUPLEmark
-> ( // apply a tuple
    TA := Contents(DATA TA); // rv of rand
    nct( (CCDE TA = RVmark) & (Isinteger (DATA TA) ) )
        -> Report 'Tuple mis-applied.' !
    TE := (Contents(DATA T)) (DATA TA);
    Stack(LVmark, & TE)
)
!
CCDE T = BASICmark // apply a basic
-> Stack( (DATA T), TA)
! Report 'Improper RATOR for GAMMA.';
gotc Lcop;

ER7: // beta -- conditional branch
TA := Pop S; // this is to be a truth value, so check
nct( (CCDE TA = RVmark) & (Iscclean(DATA TA)) )
    -> Report 'Non-boolean argument to conditional' !
TE := Pop C; // the 'false' exit
TC := Pop C; // the 'true' exit
Push( (DATA TA -> TC ! TE), C);
gotc Lcop;

ER8: // delta
LastDELTa := DATA T; // item to load into control
TA := & LastDELTa;
L8A: // copy the control into C
Push( (TA 1, TA 2), C);
TA := TA 3;
gotc Null TA -> Lcop ! L8A;

ER9: // tau -- make a k-tuple
TA, TE := nil, DATA T;
L9A: TE > 0
    -> ( TA := TA aug DATA(Pop S);
        TE := TE - 1;
        gotc L9A
    )
    !
Stack(THELEMARK, Store TA);
gotc Lcop;

ER10: // curly S
Stack(LVmark, Store(Pop S));
gotc Lcop;

ER11: // curly C
Stack(Contents(DATA(Pop S)));
gotc Lcop;

ER12: // update (:=)
TA := Pop S; // the left side
TE := Pop S; // the right side
Update(DATA TA, TE);

```

```

Stack(LVmark, C);
gctc lcop;

ER13: // semi-colon
TA := Pop S;
(CCDE TA = LVmark) & (DATA TA = 0)
-> gotc lcop
! Report 'Improper sequence element.';

ER14: // RVMAPMA
TA := DATA(Eop S); // the rator
TA := TA ( DATA(Eop S) ); // apply to first rator
DATA T = 2 -> (TA := TA(DATA(Eop S))) ! dummy;
Stack(FVmark, & TA);
gctc lcop;

ER15: // Y -- make a closure recursive
TA := Pop S; // the closure to which we are applying Y
Push( (mEQUALCOLCN, nil), C );
Push( (uCURLyC, nil), C );
Push( (mGAMMA, nil), C );
TE := Store(GUESS nil, );
Stack(LVmark, & TE);
Stack(LVmark, & TE);
Stack TA;
gctc lcop;

ER16: // equalcolon (=)
TA := Top S; // value to be stored
TE := Top S; // the place to put it
Update(DATA TE, TA); // update the memory
Stack TA; // leave an answer
gctc lcop;

L17: // BIGDELTA -- define labels
// DATA has the form (n, (L1,L1), (L2,L2), ...)
TA := DATA T;
TE := Order TA; // number of labels to define (plus 1)
L17A: TB < 2 -> gctc L17E ! dummy; // out if done
      TC := Store(EPmark, (TA TB 2, & CompCount) );
      Enter( TA TE 1, TC, CURC );
      Stack(FVmark, & CURC );
      Push( (mENV, & CURC), C );
      TB := TB - 1;
      gctc L17A;
L17F: AddDump nil;
      Push( (mDELTA, TA 1), C );
      gctc Lcce;

ER18: // gctc (body, dump)
TA := Pop S; // place to go
CCDE TA = EPmark // check validity of target
-> dummy // it's OK
! Report 'Improper raddr for goto.';
RestoreDump(DATA TA 2);
Push( (mDELTA, DATA TA 1), C );

```

```

goto Lloop;

ER19: // val
TA := Top S; // the item whose value we want
TE := Store (ENVmark, $ DumpCount); // a dump marker
EnterE(RES, TE, Cure); // put RES into the environment
Stack (ENVmark, $ Cure);
Push( (mENV, $ Cure), C );
AddDump nil;
Push( (mELTA, DATA TA 1), C );
goto Lloop;

```

```

ER20: // res
TA := Contents( DATA (Top S) ); // the dump
TE := Top S; // the result to return
RestoreDump( DATA TA );
Stack TE;
goto Lloop;

```

// This ends the Flackcard Evaluator.

// ****

// Here is the console reading loop...

```

MAIN: SA := CycleCount PALcycles;
PALcycles := CycleCount 0;
Write ( 'R: ', SA, '*n*n');

M1: SA := EMETY;
SE := Search nil;
! SE = NEWLINE -> goto M2
! SE = '!*s' -> goto M1
! dummy;
SA := SA %Conc SE;
goto M1;

M2: SA = EMPTY -> goto M1 ! dummy;
Write 'W*n';
SA = 'eval' -> goto START
! SA = 's' -> Intfrint S
! SA = 'debug' -> Debug nil
! SA = 'jdebug' -> JDebug CLEjj
! SA = 'more' -> (MoreCY nil; goto Lloop)
! SA = 'mysys' -> MYSYS nil
! SA = 'oldsys' -> OLDSYS nil
! SA = 'cycle' -> Write(CycleCount 0, ' cycles.*n')
! SA = 'lccp' -> (Lccp := CLDLccp)
! SA = 'test' -> Test 1
! SA = 'quit' -> goto Quit
! Write '?*n';

goto MAIN;

```

```
START: // begin here to do an evaluation.  
InitALL nil; // Initialize C, S, E, M, and other things.  
gctc lloop;  
  
Done: // come here at the end of the evaluation  
Write (IntCYCLES - StartCYCLES, ' interpretation cycles.*n');  
goto MAIN;  
  
Quit: Write 'Good bye.'
```

/L602/ EAL 05/14/68 0025.2 1418

```

// Art Evans' private library of PAL functions.

// Last modified on 05/12/68 at 17:11 by Evans.

def Write x = // write a tuple, without commas and parens at top
  Istuple x
  -> W(1, Crder x)
  ! Print x
where rec W(i, n) =
  n = 0 -> Print nil
  ! i > n -> dummy
  ! ( Print(x i); W(i+1, n) )

def FINISH = // call FINISH() to terminate execution
  jj ( 11().nil )

def // convert a character to an integer
Zero = 48 // ASCII value for '0'
within
CtoI x = Stoi x + Zero

def // DEBUG - written by Barkalow, typed by Evans.
Chkind x = // check the kind of x
  let y = CtoI x // convert character to numeric
  in
    y>47 & y<58 -> y-48 // digit
    ! y < 33 -> 10 // control character (sp, tab, nl)
    ! y = 59 -> 11 // semicolon
    ! y = 44 -> 11 // comma
    ! y = 40 -> 11 // left parenthesis
    ! y = 41 -> 11 // right parenthesis
    ! y = 46 -> 11 // period
    ! y = 39 -> 12 // quote
    ! -1 // default
  within

  Ch = nil // the last scanned character
  and
  Sym, Val = nil, nil // last token scanned
  within

  Nextsymbol() = // read next symbol
  ( let N, Kind = 0, 0 // two temps
  in

    SpaceLOOP: // loop to here until a non-space
    Kind := Chkind Ch; // check kind of character

```

```

Kind = 10 // space, tab or newline
-> (Ch := Readch nil; gctc SpaceICCP) // skip spaces
!
Kind = 11 // ; , ( ) .
-> (Symt := Ch; Ch := Readch nil)
!
Kind = 12 // begin a quotation
-> ( Symt, Val := 'C', '';
CHARICCP: // the loop to scan a string
    Ch := Readch nil; // read next character
    Ch = '*' // done with string
-> ( Ch := Readch nil; goto RETURN)
! Ch = '**' // process a special
-> ( Ch := Readch nil; // read next
    Ch := Ch='t' -> '*t'
    ! Ch='n' -> '*n'
    ! Ch='s' -> '*s'
    ! Ch='L' -> '*L'
    ! Ch // default
)
! dummy; // a character has been read
Val := Val &Conc Ch; // build up string
goto CHARICCP
)
! // it must be an identifier
( Symt, Val := 'C', '';
IDICCP: // build up a name or constant
Kind < 10
-> ( Val := Val &Conc Ch;
    Kind < 0
    -> (Symt := 'V')
    ! (N := 10*N + Kind);
    Ch := Readch nil;
    Kind := Chkind Ch;
    gctc IDICCP
)
!
Symt = 'C'
-> (Val := N)
!
RETURN: dummy
)
)
within

EDebug j = // finally, the function being defined...
let Lockup S = // the function that looks up identifiers
  S = 'nil' -> nil ! LookupinJ(S, j)

in
let rec // Rbasic, Rterm and Rexp are mutually recursive
( Rbasic () = // read an item
  val(let A =
    Symt = 'V' -> Lockup Val
    ! Symt = 'C' -> & Val
    ! Symt = '(' -> (Nextsym nil; Rexp nil)

```

```

        ! res nil
    in
    Nextsymbol nil;
    res A
)
and Rterm f =
  Symt = 'y' | Symt = 'c' | Symt = '('
-> Rterm( f (Rbasic nil) )
!
and Rexp () = // read an expression
let A = nil
in
bEXPLCCE: // ?
A := A aug Rterm(Rbasic nil);
SymL = ','
-> (Nextsymbol nil; goto BEXPLCCE)
!
Order A = 1
-> A 1 // don't return a 1-tuple
!
)
in

// Here execution begins for debug.
Write 'Debug entered.*n*n';
Ch := Readch nil; // start things off
LCCE: // debug's main loop
Nextsymbol nil;
Write(Rexp nil, '*n*n');
SymL = '.';
-> Write 'Program re-entered.*n'
!
goto LCCE

def Debug () = // call debug, with no preparation
  EDbug jj

def UPDATE(x, y) = // update x with y
  x := y;
  x

and GCTC x = // qcto
  gotc x

def ADebug j = // call debug, preparing SYSTEMERROR first
  let rerrE = $ SYSTEMERROR // save a copy
  in
  SYSTEMERROR := (1l x.x); // identity function
  BDebug j;
  SYSTEMERROR := rerrE

def // funny business to update SYSTEMERROR
  t x = // the value SYSTEMERROR will have is i
    let j = jj

```

```
in
#t 'System';
#Debug j;
    x // return value called with
within
#Debug =
    SYSTEMFRNG := f; // update SYSTEMFRNG
#Debug // this definition is otherwise unnecessary

def CycleCount n = // number of execution cycles
    LookupNc = n

def LookupNc = // protect it from tampering
    #LookupNo

def DoOver = // go to DoOver to restart execution
    ( l1(), L: L )
```